

Simple, Optimal, Fast, and Robust Wireless Random Medium Access Control

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Personal Research Background

- **Wireless Communications & Networking**

- ▶ Cognitive Radio
- ▶ Cooperative Communications
- ▶ OFDM/CDMA Networks
- ▶ Wireless Multimedia
- ▶ Wireless MAC

- **Network Management & Economics**

- ▶ Pricing & Revenue Management
- ▶ Service Provider Competitions
- ▶ Network Disruption Management
- ▶ Robust Network Optimization

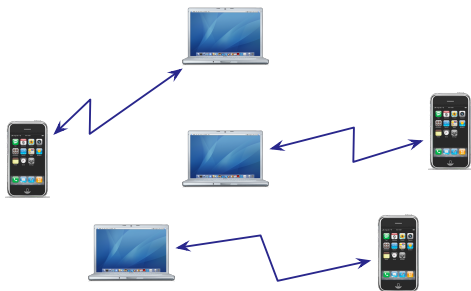
Key Methodologies

- Game Theory & Microeconomics
- Nonlinear Optimization
- Queueing & Stochastic Control

Wireless Random MAC

Ack: A. Mohsenian-Rad, M. Chiang, V. Wong

Wireless MAC Protocols



- Coordinate multiple wireless users accessing the same channel
 - ▶ Centralized: scheduling-based MAC (e.g., cellular network)
 - ▶ **Distributed: contention-based random MAC (e.g., ad hoc network)**

History of Wireless Random MAC

- Studied for 30 years
- Simplicity and practicality
- Many variations
 - ▶ Some achieved great success: Aloha, CSMA, ...
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 - ▶ Many are engineering ad hoc designs
- Two approaches:
 - ▶ **Reverse Engineering**: understand the math behind existing protocols [JSAC-07]
 - ▶ **Forward Engineering**: design better protocols [This Talk]

Forward Engineering

- How to design a better algorithm

Forward Engineering

- How to design a better algorithm
- Simple, Optimal, Fast, and Robust

Forward Engineering

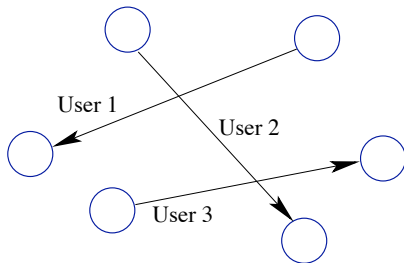
- How to design a better algorithm
- Simple, Optimal, Fast, and Robust
- Overcome performance bottlenecks of many previous algorithms

Our focus: Aloha



Aloha

(A Simple) Network Model



- A set of $\mathcal{N} = \{1, \dots, N\}$ single-hop users
- Full interference topology (relaxed later)
- Each user i
 - ▶ Contend the channel with probability $p_i \in \mathcal{P}_i = [P_i^{\min}, P_i^{\max}]$
 - ▶ Maximum data rate γ_i
 - ▶ Long term average data rate

$$r_i(\mathbf{p}) = \gamma_i p_i \prod_{j \in \mathcal{N} \setminus \{i\}} (1 - p_j)$$

Network Utility Maximization

- Each user i has an increasing and concave utility function $u_i(r_i)$

System Objective: Network Utility Maximization (NUM)

$$\max_{\mathbf{p} \in \mathcal{P}} \sum_{i \in \mathcal{N}} u_i(r_i(\mathbf{p})),$$

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- We will focus on the α -fair utility function:

$$u_i(x) = \begin{cases} (1 - \alpha)^{-1} x^{1-\alpha}, & \text{if } \alpha \in (0, 1) \cup (1, \infty), \\ \log x, & \text{if } \alpha = 1. \end{cases}$$

- ▶ $\alpha \rightarrow 0$: system throughput maximization
- ▶ $\alpha = 1$: proportional fair allocation
- ▶ $\alpha \rightarrow \infty$: max-min fairness

Previous Work

- Lee, Chiang, Canderbank 2007
- Wang, Kar 2006
- Chen, Low, Doyle 2005
- Gupta, Stolyar 2006

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- Several **performance bottlenecks**

Technical Challenges / Performance Bottlenecks

- Non-convexity: $\alpha \in (0, 1)$ is an **open** problem even centrally
- No centralized controller \Rightarrow need to be **distributed** and **asynchronous**
- Wireless lossy channels \Rightarrow messages may get **delayed** and **dropped**
- Channels can be time varying \Rightarrow demand **fast** convergence

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We will address these challenges

Key idea: Localize the Global Optimization Problem

- Intuition: each user optimizes the total network utility \Rightarrow solve NUM
- Challenge: what will be the information needed?

Local Optimization Problem

User i 's Local NUM Problem

$$\max_{p_i \in \mathcal{P}_i} \sum_{j \in \mathcal{N}} u_j(r_j(p_i, \mathbf{p}_{-i})),$$

- Objective: total network utility
- Variable: user i 's transmission probability p_i
- Parameter: other users' transmission probabilities

$$\mathbf{p}_{-i} = (p_1, \dots, p_{i-1}, p_{i+1}, \dots, p_N)$$

Optimal Solution of Local Optimization

$$p_i^*(\mathbf{p}_{-i}) = f_i(\mathbf{p}_{-i}) = \left[1 / \left(1 + \sqrt[\alpha]{v_i(\mathbf{p}_{-i})} \right) \right]_{P_i^{\min}}^{P_i^{\max}},$$

Optimal Solution of Local Optimization

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- $m_j = (1/\gamma_j)^{\alpha-1} (1/p_j - 1)^{\alpha-1}, \quad \forall j \in \mathcal{N}$

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- $m_j = (1/\gamma_j)^{\alpha-1} (1/p_j - 1)^{\alpha-1}, \quad \forall j \in \mathcal{N}$
- If each j broadcasts the message $m_j \Rightarrow$ user i can calculate $p_i^*(\mathbf{p}_{-i})$.

Local Algorithm for User i

- 1: Initialize p_i and $\mathbf{m} = (m_1, \dots, m_N)$.
- 2: **repeat**
- 3: Transmit with probability p_i .
- 4: At a **randomly** chosen time, Update

$$p_i = \left[1 / \left(1 + \sqrt[\alpha]{\gamma_i^{\alpha-1} \sum_{j \in \mathcal{N} \setminus \{i\}} m_j} \right) \right]_{p_i^{\min}}^{p_i^{\max}} .$$

- 5: At a **randomly** chosen time, update and broadcast

$$m_i = (1/\gamma_i)^{\alpha-1} (1/p_i - 1)^{\alpha-1} .$$

- 6: **until** the user decides to leave the network.

Algorithm Properties

Theorem

Under proper technical conditions and for **any** α -fair utility function:

- 1 **Uniqueness**: the algorithm has a **unique** fixed point.
- 2 **Optimality**: it is also the unique **global optimal** solution of NUM problem.
- 3 **Convergence**: the algorithm **globally** and **asynchronously** converges.
- 4 **Robustness**: convergence is **robust** to any bounded message delay/loss.

Key Maths

- ① **Uniqueness:** contraction mapping or monotonic mapping

$$\mathbf{f}(\mathbf{p}) = (f_i(\mathbf{p}_i), \forall i)$$

- ② **Optimality:** fixed point set of algorithm = KKT point set of NUM

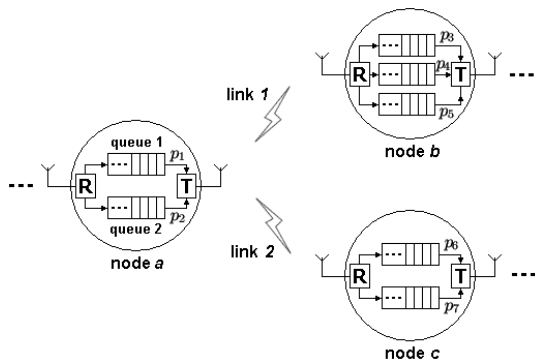
- ③ **Convergence & Robustness:**

- ▶ Synchronous convergence
- ▶ Box condition

- ④ No convexity is required

- ▶ The proposed algorithm works with enough contention level

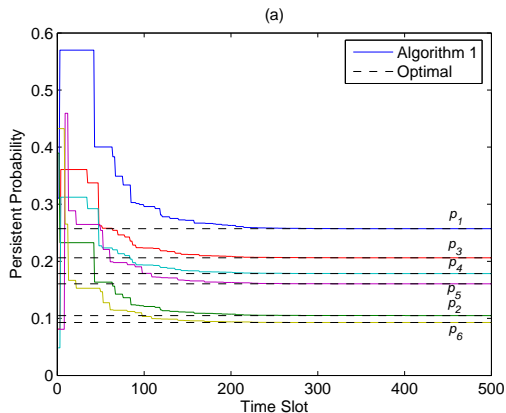
Extensions



- Can also be extended to general interference case
 - ▶ A node can have multiple outgoing links
 - ▶ A link may only interfere with a subset of other links
- All previous results go through.

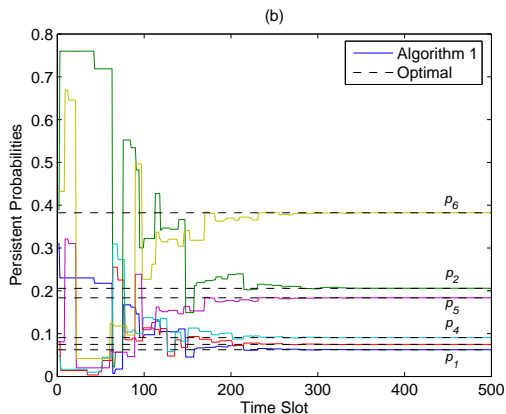
Convergence and Optimality

- 3 nodes and 6 links; $\alpha = 2$ (convex NUM)

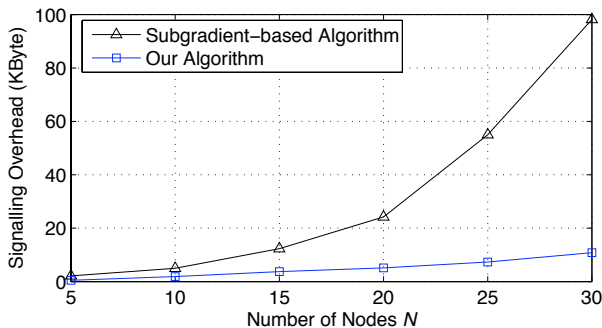


Convergence and Optimality

- 3 nodes and 6 links; $\alpha = 0.6$ (non-convex NUM)

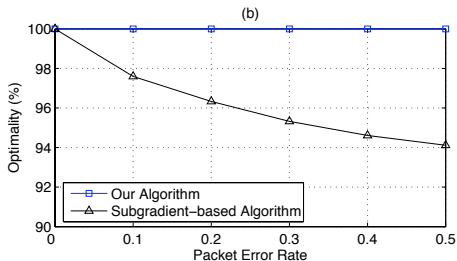
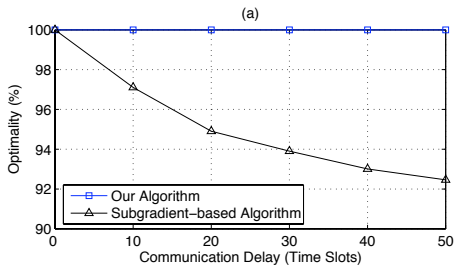


Signalling Overhead



Subgradient-based algorithm: J. Lee, M. Chiang, and R. Calderbank,
"Utility-optimal random-access control," IEEE Trans. Wireless Comm., 2007.

Impact of Delay and Message Loss



Summary

- **Topic:** Forward engineering MAC as an optimization problem
- **Algorithm:** distributed asynchronous updates with limited message passing
- **Properties:** simple, optimal, fast, and robust
- **Extension:** the same algorithm works without any explicit message passing (with limited topology)

Related Journal Publications



A. H. Mohsenian-Rad, J. Huang, M. Chiang and V.W.S. Wong, "Utility-Optimal Random Access: Reduced Complexity, Fast Convergence, and Robust Performance," *IEEE Transactions on Wireless Communications*, Feb. 2009



A. H. Mohsenian-Rad, J. Huang, M. Chiang and V.W.S. Wong, "Utility-Optimal Random Access: Optimal Performance Without Frequent Explicit Message Passing," *IEEE Transactions on Wireless Communications*, March 2009

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